CS686: Proximity Queries

Sung-Eui Yoon (윤성의)

Course URL: <u>http://sgvr.kaist.ac.kr/~sungeui/MPA</u>



Presentation Guideline: Expectations

Good summary, not full detail, of the paper

- Just 20 min for the talk + 3 min for quiz
- Talk about motivations of the work
- Give a broad background on the related work
- Explain main idea and results of the paper
- Discuss strengths and weaknesses of the method



High-Level Ideas

Deliver most important ideas and results

- Do not talk about minor details
- Give enough background instead

- Spend most time to figure out the most important things and prepare good slides for them
 - If possible, re-use existing slides/videos with acknowledgement



Overall Structure

• Prepare an overview slide

Talk about most important things and connect them well



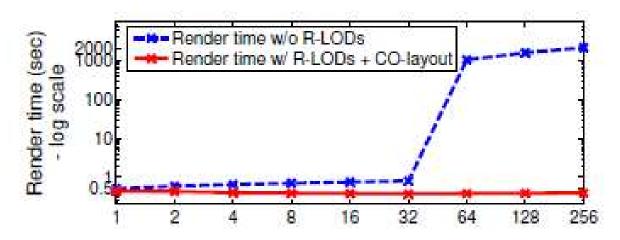
Be Honest

- Do not skip important ideas that you don't know
 - Explain as much as you know and mention that you don't understand some parts
- If you get questions you don't know good answers, just say it
- In the end, you need to explain them before the semester ends



Result Presentation

Give full experiment settings and present data with the related information



- After showing the data, give a message that we can pull of the data
- Show images/videos, if there are



Prepare a Quiz

- Give two simple questions to draw attentions
 - Ask a keyword
 - Simple true or false questions
 - Multiple choice questions
- Grade them in the scale of 0 and 10, and send the score to TA



Audience feedback form

Date: Talk title: Speaker:

A. Was the talk well organized and well prepared?5: Excellent 4: good 3: okay 2: less than average 1: poor

B. Was the talk comprehensible? How well were important concepts covered?5: Excellent 4: good 3: okay 2: less than average 1: poor

Any comments to the speaker



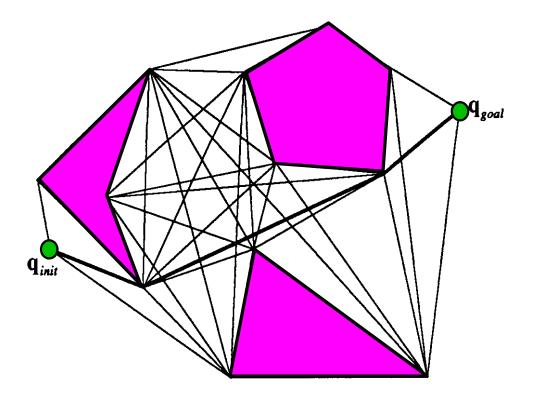
Class Objectives

- Understand collision detection and distance computation
 - Bounding volume hierarchies
- Handle point clouds



Two geometric primitives in configuration space

- CLEAR(q)
 Is configuration q collision
 free or not?
- LINK(q, q')
 Is the straight-line path
 between q and q'
 collision-free?





Problem

- Input: two objects A and B
- Output:
 - Distance computation: compute the distance (in the workspace) between *A* and *B*

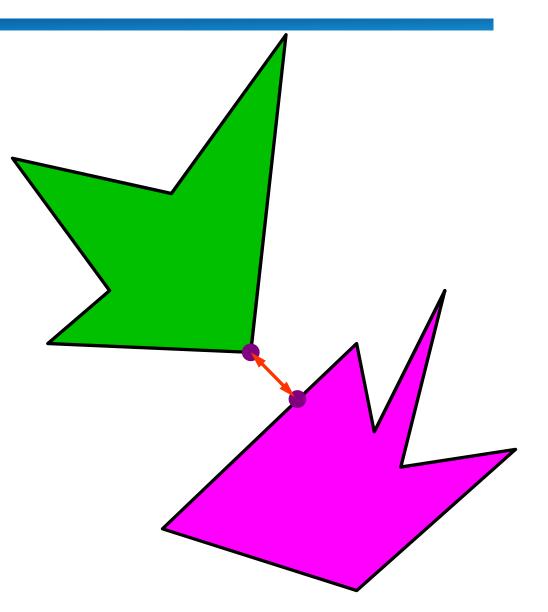
OR

• Collision detection: determine whether A and B collide or not



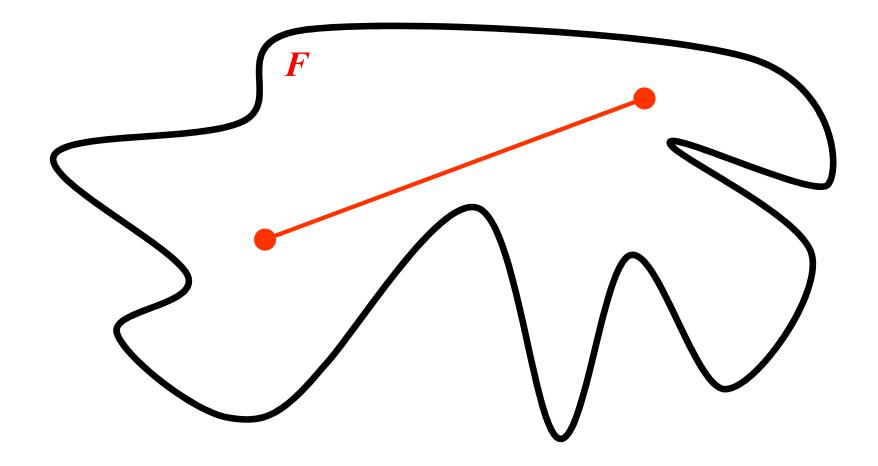
Collision detection vs. distance computation

- The distance between two objects (in the workspace) is the distance between the two closest points on the respective objects
- Collision if and only if distance = 0



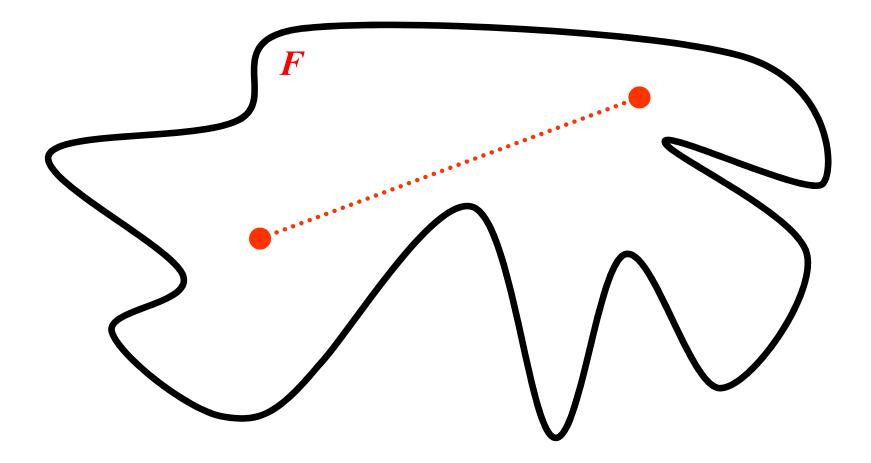


Collision detection does not allow us to check for free path rigorously





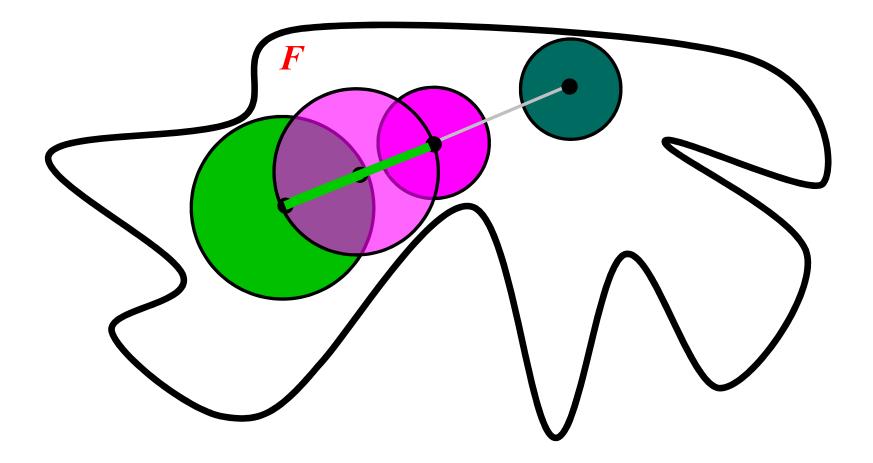
Collision detection does not allow us to check for free path rigorously



Discrete collision checks



Use distance to check for free path rigorously





Use distance to check for free path rigorously

Link(q0, q1)

1: if $q0 \in N(q1)$ or $q1 \in N(q0)$

2: then

- 3: return TRUE.
- 4: else
- 5: q' = (q0+q1)/2.
- 6: if q' is in collision

7: then

8: return FALSE

9: else

10: return Link(q0, q') & Link(q1,q') KAIST

Applications

- Robotics
 - Collision avoidance
 - Path planning
- Graphics & virtual environment simulation
- Haptics
 - Collision detection
 - Force proportional to distance



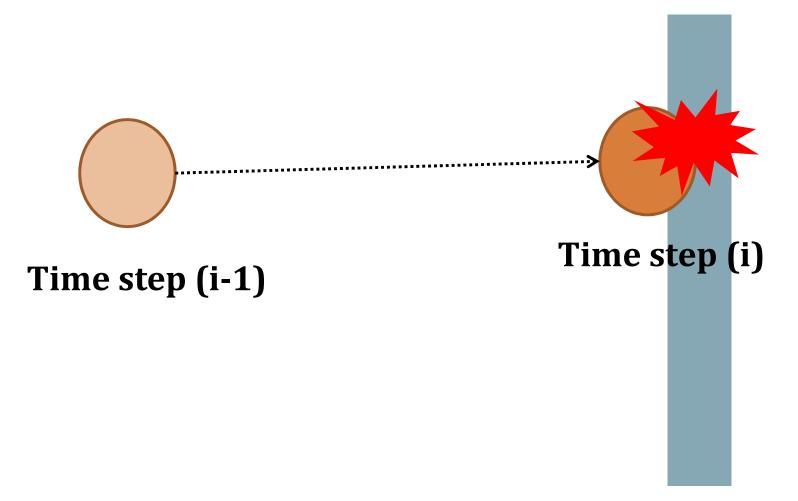


Collision Detection

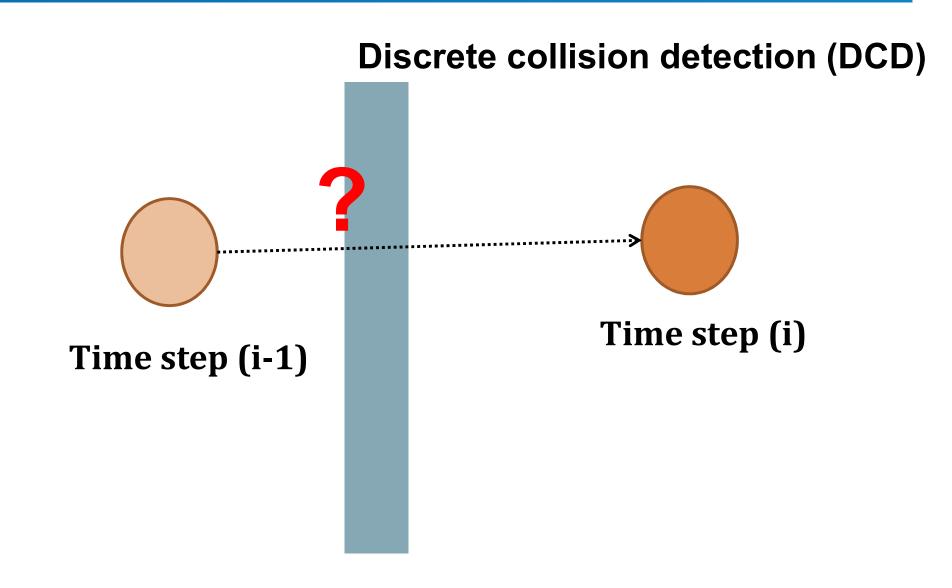
- Discrete collision detection
- Continuous collision detection



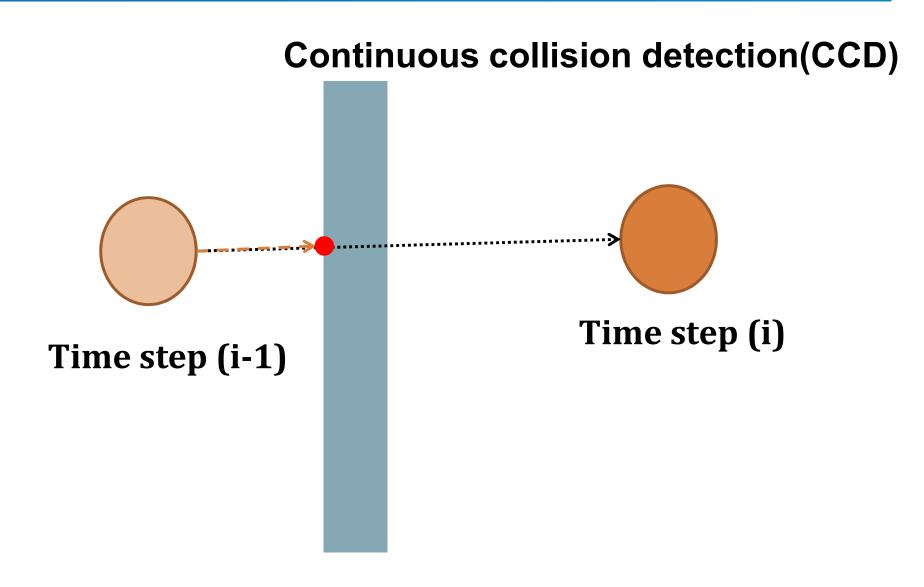
Discrete collision detection (DCD)













	Continuous CD	Discrete CD
Accuracy	Accurate	May miss some collisions
Computation time	Slow	Fast



Collision Detection

- Discrete collision detection
- Continuous collision detection
- These are typically accelerated by bounding volume hierarchices (BVHs)



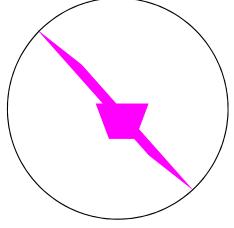
Bounding Volumes

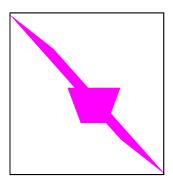
Sphere [Whitted80]

- Cheap to compute
- Cheap test
- Potentially very bad fit

Axis-aligned bounding box

- Very cheap to compute
- Cheap test
- Tighter than sphere



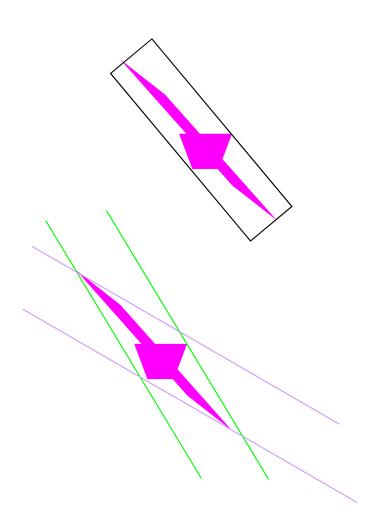




Bounding Volumes

Oriented bounding box

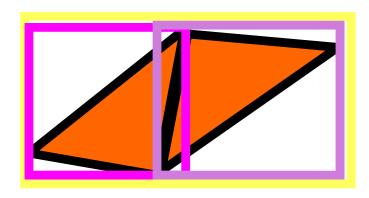
- Fairly cheap to compute
- Fairly cheap test
- Generally fairly tight
- Slabs / K-dops
 - More expensive to compute
 - Fairly cheap test
 - Can be tighter than OBB

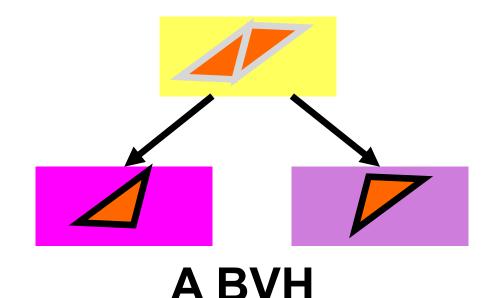




Bounding Volume Hierarchies (BVHs)

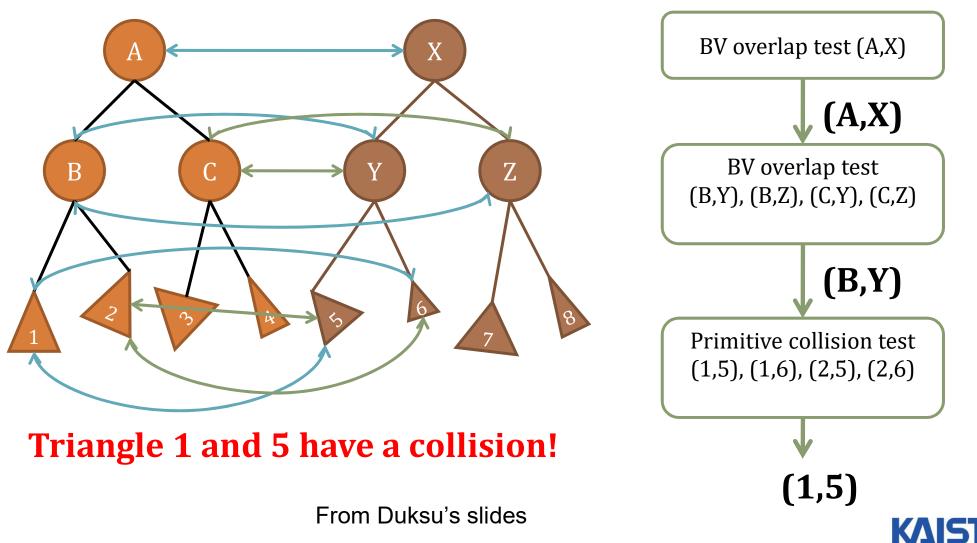
- Organize bounding volumes recursively as a tree
- Construct BVHs in a top-down manner
 - Use median-based partitioning or other advanced partitioning methods







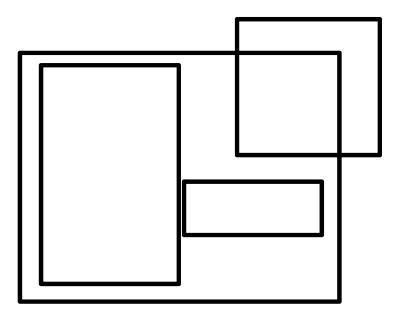
Collision Detection with BVHs



BVH Traversal

 Traverse BVHs with depth-first or breadthfirst

• Refine a BV node first that has a bigger BV





Continuous Collision Detection

- BVHs are also widely used
- Models a continuous motion for a primitive, whose positions are defined at discrete time steps
 - E.g., linear interpolation



Test-Of-Time 2006 Award



RT-DEFORM: Interactive Ray Tracing of Dynamic

Scenes using BVHs

Christian Lauterbach, Sung-eui Yoon, David Tuft,

Dinesh Manocha

IEEE Interactive Ray Tracing, 2006



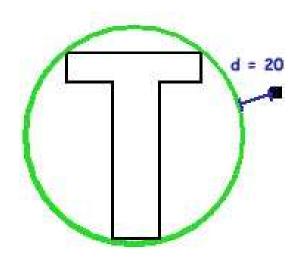
Computing distances

Depth-first search on the binary tree

- Keep an updated minimum distance
- Depth-first \rightarrow more pruning in search
- Prune search on branches that won't reduce minimum distance
- Once leaf node is reached, examine underlying convex polygon for exact distance



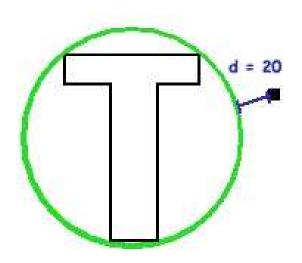
Set initial distance value to infinity



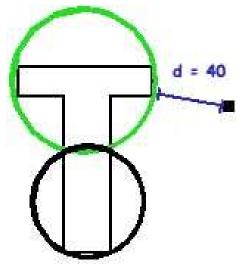
Start at the root node. 20 < infinity, so continue searching



Set initial distance value to infinity



Start at the root node. 20 < infinity, so continue searching.

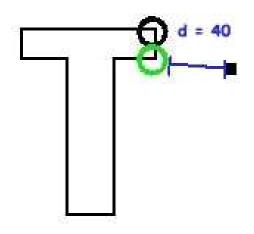


40 < infinity, so continue searching recursively.

Choose the nearest of the two child spheres to search first



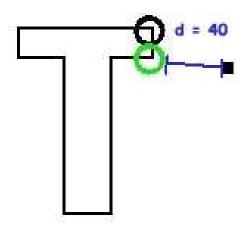
• Eventually reach a leaf node

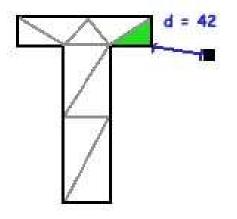


40 < infinity; examine the polygon to which the leaf node is attached.



Eventually reach a leaf node

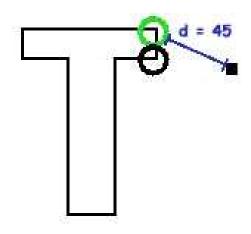




40 < infinity; examine the polygon to which the leaf node is attached.

Call algorithm to find exact distance to the polygon. Replace infinity with new minimum distance (42 in this case).

Continue depth-first search

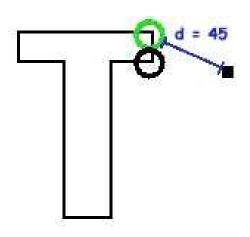


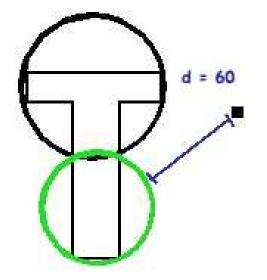
45 > 42; don't search this branch any further



Simple example

• Continue depth-first search





45 > 42; don't search this branch any further 60 > 42; we can prune this half of our tree from the search



Running time: build the tree

- Roughly balanced binary tree
- Expected time $O(n \log n)$
 - Time to generate node v is proportional to the number of leaf nodes descended from v.
- Tree is built only once and can often be pre-computed.



Running time: search the tree

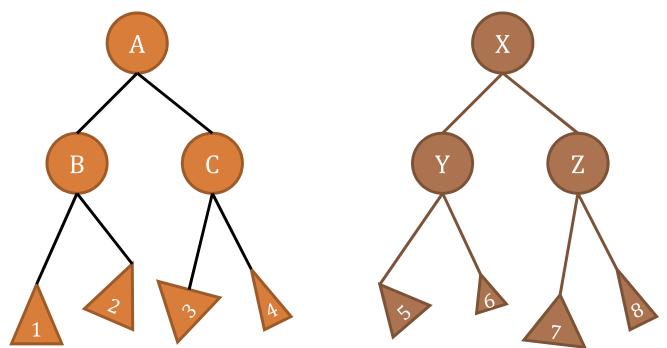
Full search

- O(n) time to traverse the tree, where n = number of leaf nodes
- Plus time to compute distance to each polygon in the underlying model
- The algorithm allows a pruned search:
 - Worst case as above; occurs when objects are close together
 - Best case: O(log n) + a single polygon calculation
 - Average case ranges between the two





- If second object is not a single point, then search & compare 2 trees
 - Use two BVHs and perform the BVH traversal

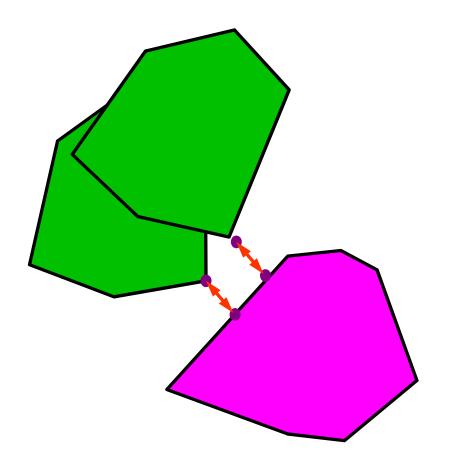




Tracking the closest pair

V-Clip: Fast and Robust Polyhedral Collision Detection, B. Mirtich, 1997

Utilize motion coherence





Sensor-based Path Planning

Navigation using 3D depth sensor

Real-Time Navigation in 3D Environments Based on Depth Camera Data

Daniel Maier Armin Hornung Maren Bennewitz

Humanoid Robots Laboratory, University of Freiburg



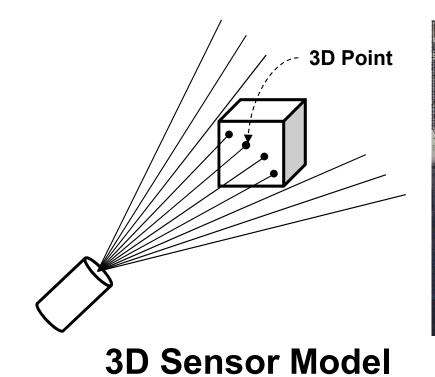
Maier, Daniel, et al.

Modified from YongSun's slides



3D Sensor & Point Cloud Data

- 3D sensor generates excessive amount of points with some noise periodically
 - 300K points / 30FPS with Kinect

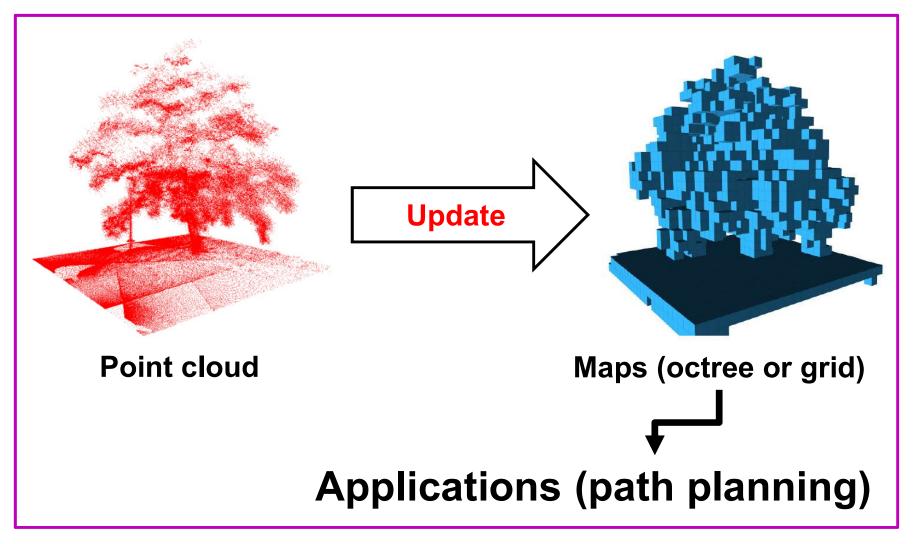




Point Cloud Data

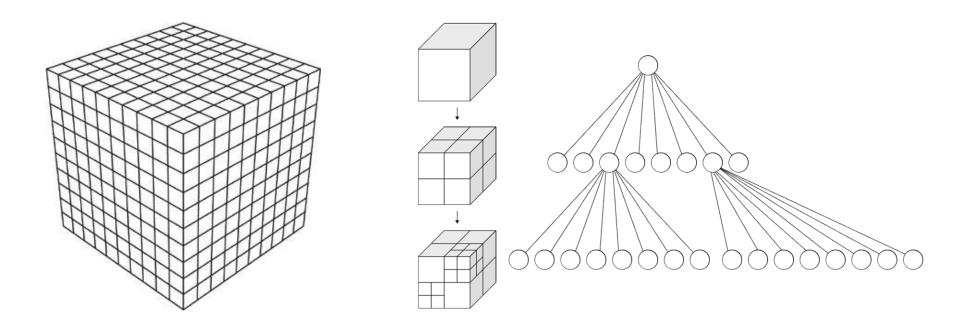


General Flow of Using Point Clouds





Map Representations



3D Grid Map

Octree Data Structure



Occupancy Map Representation

OctoMap [Wurm et al., ICRA, 2010]

• Encode an occupancy probability of cell *n* given measurement *z*_{1:*t*}

$$L(n \mid z_{1:t}) = L(n \mid z_{1:t-1}) + L(n \mid z_t)$$

Occupancy probability of the
cell *n* at time step *t* - 1
New sensor measurement *z*_t

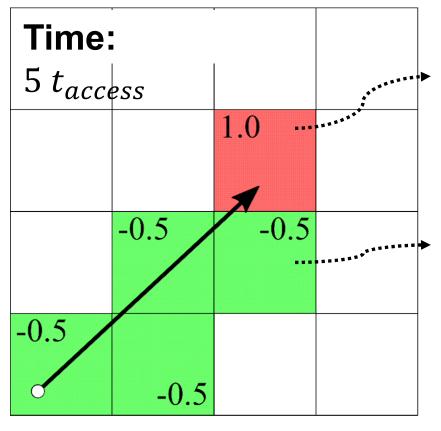
to be updated at time step t $L(n \mid z_t) = \begin{cases} l_{occ} & occupied state \\ l_{free} & free state \end{cases}$



Update Method

Traverse and update cells

 Bresenham algorithm [Amanatides et al., *Eurographics*, 1987]



Updated cell to occupied state $L(n \mid z_t) = l_{occ} = 1.0$

Updated cell to free state $L(n \mid z_t) = l_{free} = -0.5$

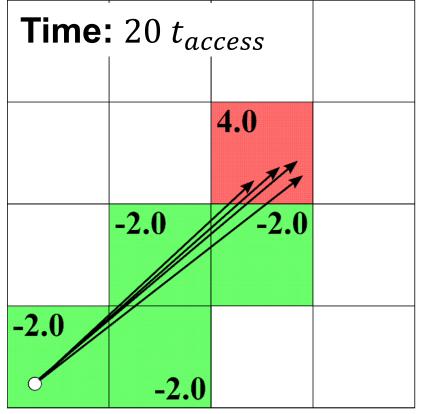
t_{access}: time to update a cell



Update Method

Traverse and update cells

- Bresenham algorithm [Amanatides et al., Eurographics, 1987]
- Can be very slow, with many points



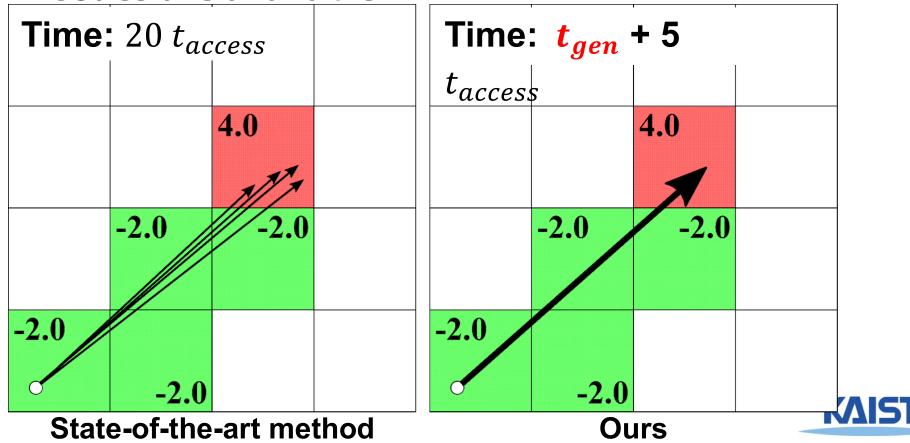
- Visit the same cells multiple times for multiple rays



Super Rays [Kwon et al., ICRA16]

Benefits of our approach

- Faster performance with the same representation accuracy
- Codes are available



Class Objectives were:

- Understand collision detection and distance computation
 - Bounding volume hierarchies
- Handle point clouds



Next Time...

Probabilistic Roadmaps



Homework

- Submit summaries of 2 ICRA/IROS/RSS/CoRL/TRO/IJRR papers
- Go over the next lecture slides
- Come up with 3 questions before the midterm exam

