## CS686: <br> Proximity Queries

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Course URL:
http://sglab.kaist.ac.kr/~sungeui/MPA

## KAIST

## Presentation Guideline: Expectations

- Good summary, not full detail, of the paper
- Just 15min ~ 20 min
- Talk about motivations of the work
- Give a broad background on the related work
- Explain main idea and results of the paper
- Discuss strengths and weaknesses of the method


## High-Level Ideas

- Deliver most important ideas and results
- Do not talk about minor details
- Give enough background instead
- Spend most time to figure out the most important things and prepare good slides for them
- If possible, re-use existing slides/ videos with ack.


## Overall Structure

- Prepare an overview slide
- Talk about most important things and connect them well


## Be Honest

- Do not skip important ideas that you don't know
- Explain as much as you know and mention that you don't understand some parts
- If you get questions you don't know good answers, just say it
- In the end, you need to explain them before the semester ends


## Result Presentation

- Give full experiment settings and present data with the related information

- After showing the data, give a message that we can pull of the data
- Show images/ videos, if there are


## Prepare a Quiz

- Give two simple questions to draw attentions
- Ask a keyword
- Simple true or false questions
- Multiple choice questions
- Grade them in the scale of 0 and 10, and send the score to TA


## Audience feedback form

## Date:

Talk title:
Speaker:
A. Was the talk well organized and well prepared?

5: Excellent 4: good 3: okay 2: less than average 1: poor
B. Was the talk comprehensible? How well were important concepts covered?

5: Excellent 4: good 3: okay 2: less than average 1: poor
Any comments to the speaker

## Class Objectives

- Understand collision detection and distance computation
- Bounding volume hierarchies
- Handle point clouds


## Two geometric primitives in configuration space

- Clear(q) Is configuration q collision free or not?
- LiNK(q, q')

Is the straight-line path between $q$ and $q$ ' collision-free?


## Problem

- Input: two objects $A$ and $B$
- Output:
- Distance computation: compute the distance (in the workspace) between $A$ and $B$ OR
- Collision detection: determine whether $A$ and $B$ collide or not


## Collision detection vs. distance computation

- The distance between two objects (in the workspace) is the distance between the two closest points on the respective objects
- Collision if and only if distance $=0$



## Collision detection does not allow us to check for free path rigorously



## Collision detection does not allow us to check for free path rigorously



Discrete collision checks

## Use distance to check for free path rigorously



## Use distance to check for free path rigorously

Link(q0, q1)
1: if $q 0 \in N(q 1)$ or $q 1 \in N(q 0)$
2: then
3: return TRUE.
4: else
5: $\quad q^{\prime}=(q 0+q 1) / 2$.
6: if q' is in collision
7: then
8: return FALSE
9: else

```
16 10: return Link(q0, q') && Link(q1,q') KAIST
```


## Applications

- Robotics
- Collision avoidance
- Path planning
- Graphics \& virtual environment simulation
- Haptics
- Collision detection
- Force proportional to distance



## Collision Detection

- Discrete collision detection
- Continuous collision detection


## Discrete VS Continuous

## Discrete collision detection (DCD)



## Discrete VS Continuous

Discrete collision detection (DCD)

Time step (i-1)
Time step (i)

## Discrete VS Continuous

## Continuous collision detection(CCD)



## Discrete VS Continuous

|  | Continuous CD | Discrete CD |
| :---: | :---: | :---: |
| Accuracy | Accurate | May miss some collisions |
| Computation time | Slow | Fast |

## Collision Detection

- Discrete collision detection
- Continuous collision detection
- These are typically accelerated by bounding volume hierarchices (BVHs)


## Bounding Volumes

- Sphere [Whitted80]
- Cheap to compute
- Cheap test
- Potentially very bad fit
- Axis-aligned bounding box
- Very cheap to compute
- Cheap test
- Tighter than sphere



## Bounding Volumes

- Oriented bounding box
- Fairly cheap to compute
- Fairly cheap test
- Generally fairly tight
- Slabs / K-dops
- More expensive to compute
- Fairly cheap test
- Can be tighter than OBB


## Bounding Volume Hierarchies (BVHs)

- Organize bounding volumes recursively as a tree
- Construct BVHs in a top-down manner
- Use median-based partitioning or other advanced partitioning methods



A BVH

## Collision Detection with BVHs



Triangle 1 and 5 have a collision!

$(1,5)$

From Duksu's slides

## BVH Traversal

- Traverse BVHs with depth-first or breadthfirst
- Refine a BV node first that has a bigger BV



## Continuous Collision Detection

- BVHs are also widely used
- Models a continuous motion for a primitive, whose positions are defined at discrete time steps
- E.g., linear interpolation


## Test-Of-Time 2006 Award

## High-Performance Graphics 2015

Los Angeles, August 7-9, 2015

RT-DEFORM: I nteractive Ray Tracing of Dynamic Scenes using BVHs
Christian Lauterbach, Sung-eui Yoon, David Tuft, Dinesh Manocha

IEEE Interactive Ray Tracing, 2006

## Computing distances

- Depth-first search on the binary tree
- Keep an updated minimum distance
- Depth-first $\rightarrow$ more pruning in search
- Prune search on branches that won't reduce minimum distance
- Once leaf node is reached, examine underlying convex polygon for exact distance


## Simple example

- Set initial distance value to infinity


Start at the root node. 20 < infinity, so continue searching

## Simple example

- Set initial distance value to infinity


Start at the root node. 20 < infinity, so continue searching.

$40<$ infinity, so continue searching recursively.

- Choose the nearest of the two child spheres to search first


## Simple example

- Eventually reach a leaf node

$40<$ infinity; examine the polygon to which the leaf node is attached.


## Simple example

- Eventually reach a leaf node


40 < infinity; examine the polygon to which the leaf node is attached.


Call algorithm to find exact distance to the polygon.
Replace infinity with new minimum distance (42 in this case).

## Simple example

- Continue depth-first search

$45>42$; don't search this
branch any further


## Simple example

- Continue depth-first search

$45>42$; don't search this branch any further

$60>42$; we can prune this half of our tree from the search


## Running time: build the tree

- Roughly balanced binary tree
- Expected time $O(n \log n)$
- Time to generate node $v$ is proportional to the number of leaf nodes descended from $v$.
- Tree is built only once and can often be pre-computed.


## Running time: search the tree

- Full search
- $O(n)$ time to traverse the tree, where $n=$ number of leaf nodes
- Plus time to compute distance to each polygon in the underlying model
- The algorithm allows a pruned search:
- Worst case as above; occurs when objects are close together
- Best case: $O(\log n)+$ a single polygon calculation
- Average case ranges between the two


## General case

- If second object is not a single point, then search \& compare 2 trees
- Use two BVHs and perform the BVH traversal



## Tracking the closest pair

- V-Clip; Fast and Robust Polyhedral Collision Detection, B. Mirtich, 1997
- Utilize motion coherence



## Sensor-based Path Planning

- Navigation using 3D depth sensor

> Real-Time Navigation in 3D Environments Based on Depth Camera Data

Daniel Maier Armin Hornung Maren Bennewitz

Humanoid Robots Laboratory, University of Freiburg


Maier, Daniel, et al.

## 3D Sensor \& Point Cloud Data

- 3D sensor generates excessive amount of points with some noise periodically
- 300K points / 30FPS with Kinect


3D Sensor Model


Point Cloud Data

## General Flow of Using Point Clouds



## Map Representations



3D Grid Map


Octree Data Structure

## Occupancy Map Representation

- OctoMap [Wurm et al., /CRA, 2010]
- Encode an occupancy probability of cell $n$ given measurement $z_{1: t}$

$$
L\left(n \mid z_{1: t}\right)=L\left(n \mid z_{1: t-1}\right)+L\left(n \mid z_{t}\right)
$$

Occupancy probability of the cell $n$ at time step $t-1$

New sensor measurement $z_{t}$ to be updated at time step $t$

$$
L\left(n \mid z_{t}\right)= \begin{cases}l_{\text {occ }} & \text { occupied state } \\ l_{\text {free }} & \text { free state }\end{cases}
$$

## Update Method

- Traverse and update cells
- Bresenham algorithm [Amanatides et al., Eurographics, 1987]


Updated cell to occupied state
$L\left(n \mid z_{t}\right)=l_{o c c}=1.0$

Updated cell to free state
$L\left(n \mid z_{t}\right)=l_{\text {free }}=-0.5$
$t_{\text {access }}$ : time to update a cell

## Update Method

- Traverse and update cells
- Bresenham algorithm [Amanatides et al., Eurographics, 1987]
- Can be very slow, with many points

- Visit the same cells multiple times for multiple rays
$t_{\text {access }}$ : time to update a cell


## Super Rays [Kwon et al., ICRA16]

- Benefits of our approach
- Faster performance with the same representation accuracy
- Codes are available


State-of-the-art method


## Class Objectives were:

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## Next Time...

- Probabilistic Roadmaps

